# DEPARTMENT OF COMPUTER SCIENCE \& ENGINEERING 

Year / Sem : II / II

Sub. Code \& Subject : 18CSE225 - Formal Languages and Automata Theory
Prepared by
: Dr. D. Jagadeesan, Professor / CSE

## QUESTION BANK

## Unit - I

| $\begin{aligned} & \text { Sl. } \\ & \text { No } \\ & \hline \end{aligned}$ | Questions | CO | PO | BT |
| :---: | :---: | :---: | :---: | :---: |
| Part - A |  |  |  |  |
| 1 | Write down the operations on set. | 1 | 1 | 1 |
| 2 | List any three applications of Automata Theory. | 1 | 1 | 1 |
| 3 | Define Finite Automation. | 1 | 1 | 1 |
| 4 | Define Deterministic Finite Automation. | 1 | 1 | 1 |
| 5 | Define Non-Deterministic Finite Automation. | 1 | 1 | 1 |
| 6 | Define NFA with $\varepsilon$ transition. | 1 | 1 | 1 |
| 7 | Design FA which accepts odd number of 1's and any number of 0's. | 1 | 2,3 | 6 |
| 8 | Design FA to check whether given unary number is divisible by three. | 1 | 2,3 | 6 |
| 9 | Design FA to check whether given binary number is divisible by three. | 1 | 2,3 | 6 |
| 10 | Design FA to accept the string that always ends with 00. | 1 | 2,3 | 6 |
| 11 | Obtain the $\varepsilon$ closure of states $q 0$ and q 1 in the following NFA with $\varepsilon$ transition. | 1 | 2 | 5 |
| 12 | Obtain $\varepsilon$ closure of each state in the following NFA with $\varepsilon$ move. | 1 | 2 | 5 |
| 13 | Explain a transition diagram. | 1 | 1 | 2 |
| 14 | Explain a transition table. | 1 | 1 | 2 |
| 15 | Explain the transition function. | 1 | 1 | 2 |
| 16 | Differentiate DFA and NFA? | 1 | 2 | 2 |
| 17 | Write notes on Moore Machine. | 1 | 1 | 6 |
| 18 | Write the formal definition of Moore Machine. | 1 | 1 | 6 |
| 19 | Short notes on Mealy Machine. | 1 | 1 | 1 |
| 20 | Write the formal definition of Mealy Machine. | 1 | 1 | 6 |
| 21 | Compare the Mealy and Moore Model? | 1 | 2 | 5 |
| Part - B |  |  |  |  |
| 22 | Design FA to accept the string that always ends with 00. | 1 | 2,3 | 6 |
| 23 | Design FA to check whether given binary number is divisible by three. | 1 | 2,3 | 6 |
| 24 | Show that "For every NFA, there exists a DFA which simulates the behavior of NFA. If L is the set accepted by NFA, then there exists a DFA which also accepts L". | 1 | 2 | 1 |
| 25 | Show that "If L is accepted by NFA with $\varepsilon$-moves, then there exists L which is accepted by NFA without $\varepsilon$-moves. | 1 | 2 | 1 |
| 26 | Construct DFA equivalent to the given NFA | 1 | 2,3 | 6 |

\begin{tabular}{|c|c|c|c|c|c|c|c|c|}
\hline 27 \& \multicolumn{5}{|l|}{Let \(\mathrm{M}=(\{\mathrm{q} 0, \mathrm{q} 1\},\{0,1\}, \delta, \mathrm{q} 0,\{\mathrm{q} 1\})\) be NFA. Where \(\delta(\mathrm{q} 0,0)=\{\mathrm{q} 0\), \(\mathrm{q} 1\}, \delta(\mathrm{q} 0,1)=\{\mathrm{q} 1\}, \delta(\mathrm{q} 1,0)=\{\phi\}, \delta(\mathrm{q} 1,1)=\{\mathrm{q} 0, \mathrm{q} 1\}\). Construct its equivalent DFA.} \& 1 \& 2,3 \& 6 \\
\hline 28 \& \multicolumn{5}{|l|}{\begin{tabular}{l}
Let \(\mathrm{M}=(\{\mathrm{q} 0, \mathrm{q} 1, \mathrm{q} 2, \mathrm{q} 3\},\{0,1\}, \delta, \mathrm{q} 0,\{\mathrm{q} 2, \mathrm{q} 3\})\) be \(\varepsilon\)-NFA. \\
Where \(\delta(\mathrm{q} 0,0)=\{\mathrm{q} 0, \mathrm{q} 1\}, \delta(\mathrm{q} 0,1)=\{\mathrm{q} 1\}, \delta(\mathrm{q} 1,0)=\{\mathrm{q} 2, \mathrm{q} 3\}, \delta(\mathrm{q} 1\), \\
\(\varepsilon)=\{\mathrm{q} 1\}, \delta(\mathrm{q} 1,1)=\{\mathrm{q} 0, \mathrm{q} 1\}, \delta(\mathrm{q} 2,0)=\{\mathrm{q} 2\}, \delta(\mathrm{q} 2, \varepsilon)=\{\mathrm{q} 3\}, \delta(\mathrm{q} 2\), \\
\(1)=\{q 0, q 3\},, \delta(q 3,0)=\{q 3\}, \delta(q 3,1)=\{q 2, q 3\}, \delta(q 3, \varepsilon)=\{q 0\}\). \\
Construct its equivalent DFA.
\end{tabular}} \& 1 \& 2,3 \& 6 \\
\hline \multirow{4}{*}{29} \& \multicolumn{5}{|l|}{Consider the following \(\varepsilon\)-NFA. Compute the \(\varepsilon\)-closure of each state and find it's equivalent DFA.} \& \multirow{4}{*}{1} \& \multirow{4}{*}{2,3} \& \multirow{4}{*}{5} \\
\hline \& \& \(\Phi\) \& \{p\} \(\{\mathrm{q}\}\) \& \(\Phi\) \& \& \& \& \\
\hline \& \& q \(\quad\{\mathrm{p}\}\) \& \{q\} \(\{\mathrm{r}\}\) \& \(\Phi\) \& \& \& \& \\
\hline \& \& *r \({ }^{\text {r }}\) [q\} \& \{r\} \(\quad\) ¢ \& \{p\} \& \& \& \& \\
\hline 30 \& \multicolumn{5}{|l|}{Convert a NFA which accepts the string ends with 01 to a DFA.} \& 1 \& 2,3 \& 5 \\
\hline 31 \& \multicolumn{5}{|l|}{Consider the Moore machine described by the transition diagram given below. To construct a Mealy machine, which is equivalent to moore machine} \& 1

1 \& 2,3 \& 5 <br>
\hline \multirow{7}{*}{32} \& \multicolumn{5}{|l|}{Consider the Mealy machine described by the transition table given below. To construct a Moore machine, which is equivalent to mealy machine?} \& \multirow{7}{*}{1} \& \multirow{7}{*}{2,3} \& \multirow{7}{*}{5} <br>
\hline \& \multirow{2}{*}{Present State} \& \multicolumn{2}{|l|}{input $=0$} \& \multicolumn{2}{|l|}{input $=1$} \& \& \& <br>
\hline \& \& Next State \& Output \& Next State \& Output \& \& \& <br>
\hline \& $\rightarrow \mathrm{q} 1$ \& q3 \& 0 \& q2 \& 0 \& \& \& <br>
\hline \& q2 \& q1 \& 1 \& q4 \& 0 \& \& \& <br>
\hline \& q3 \& q2 \& 1 \& q1 \& 1 \& \& \& <br>
\hline \& q4 \& q4 \& 1 \& q3 \& 0 \& \& \& <br>
\hline
\end{tabular}

## Unit - II

| Sl. <br> No | Questions | CO | PO | BT |
| :---: | :---: | :---: | :---: | :---: |
| Part - A |  |  |  |  |
| 1 | State regular expression. | 2 | 1 | 1 |
| 2 | How the kleen's closure of L can be denoted? | 2 | 1,2 | 4 |
| 3 | How do you represent positive closure of L? | 2 | 1,2 | 4 |
| 4 | Write the regular expression for the language accepting all combinations of a's over the set $\sum=\{a\}$. | 2 | 2,3 | 6 |
| 5 | Write regular expression for the language accepting the strings which are starting with 1 and ending with 0 , over the set $\sum=\{0,1\}$. | 2 | 2,3 | 6 |
| 6 | Show that $\left(0^{*} 1^{*}\right)^{*}=(0+1)^{*}$. | 2 | 2 | 2 |
| 7 | Show that ( $\mathrm{r}+\mathrm{s})^{*} \neq \mathrm{r}^{*}+\mathrm{s}^{*}$. | 2 | 2 | 2 |
| 8 | If $\mathrm{L}=\{$ The language starting and ending with ' $a$ ' and having any combinations of $b$ 's in between, that what is r ? | 2 | 2,3 | 4 |
| 9 | Give regular expression for $\mathrm{L}=\mathrm{L} 1 \cap \mathrm{~L} 2$ over alphabet $\{\mathrm{a}, \mathrm{b}\}$ where $\mathrm{L} 1=$ all strings of even length <br> $\mathrm{L} 2=$ all strings starting with ' b '. | 2 | 2,3 | 2 |
| 10 | Explain the application of the pumping lemma. | 2 |  | 3 |
| 11 | Describe the following by regular expression <br> a. $\mathrm{L} 1=$ the set of all strings of 0 's and 1 's ending in 00 . <br> b. L2 $=$ the set of all strings of 0 's and 1 's beginning with 0 and ending with 1 . | 2 | 2,3 | 1 |
| 12 | Show that ( $\left.\mathrm{r}^{*}\right)^{*}=\mathrm{r}^{*}$ for a regular expression r . | 2 | 2 | 2 |
| 13 | Write down the closure properties of regular language. | 2 | 3 | 6 |
| 14 | What is pumping lemma? | 2 | 2 | 4 |
| 15 | State Arden's theorem. | 2 | 1 | 1 |
| 16 | What is dead state? | 2 | 2 | 4 |
| Part-B |  |  |  |  |
| 17 | Show that ' $r$ ' be a regular expression, the there exists an NFA with $\varepsilon$ transitions that accepts L\{r\}. | 2 | 2 | 2 |
| 18 | Construct the NFA with $\varepsilon$ for the regular expression using Thomson construction method. <br> a. $0(0+1) * 100$ <br> b. $a(a+b) * b$ | 2 | 2,3 | 6 |
| 19 | Obtain the equivalent DFA from the following regular expressions <br> a. $(a+b) * a b b$ <br> b. $(00+11)^{*}(0+1)^{*}$ | 2 | 2,3 | 5 |
| 20 | Show that the following languages are not regular using pumping lemma <br> a. $L=\left\{0^{i} 1^{i} ; i>=1\right\}$ <br> b. $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{p}} ; \mathrm{p}\right.$ is prime $\}$ | 2 | 2,3 | 2 |
| 21 | Find the regular expression for the set of all strings denotes by $\left(\mathrm{R}_{13}\right)^{2}$ from the deterministic finite automata given below | 2 | 2,3 | 5 |



Unit - III

| $\begin{array}{\|l} \hline \text { Sl. } \\ \text { No } \end{array}$ | Questions | CO | PO | BT |
| :---: | :---: | :---: | :---: | :---: |
| Part - A |  |  |  |  |
| 1 | Obtain the Right Linear Grammar from the given Left Linear Grammar | 3 | 1 | 5 |
| 2 | Let $\mathrm{G}=(\{\mathrm{S}, \mathrm{C}\},\{\mathrm{a}, \mathrm{b}\}, \mathrm{P}, \mathrm{S}\}$ where P consists of $\mathrm{S} \rightarrow \mathrm{aCa}, \mathrm{C} \rightarrow \mathrm{aCa}$, Find L(G))? | 3 | 2 | 5 |
| 3 | Consider G whose productions are $\mathrm{S} \rightarrow \mathrm{aAS} / \mathrm{a}, \mathrm{A} \rightarrow \mathrm{SbA} / \mathrm{SS} / \mathrm{ba}$, show that $S \rightarrow$ aabbaa and construct a derivation tree. | 3 | 2 | 2 |
| 4 | Find L(G) where G $=(\{\mathrm{S}\},\{0,1\},\{\mathrm{S} \rightarrow 0 \mathrm{~S} 1, \mathrm{~s} \rightarrow \varepsilon\}, \mathrm{S})$ | 3 | 2 | 5 |
| 5 | Construct a CFL from the given grammar $\mathrm{S} \rightarrow \mathrm{aaA}, \mathrm{A} \rightarrow \mathrm{S} / \mathrm{a}$ | 3 | 2 | 6 |
| 6 | Define a derivation tree for CFG. | 3 | 1 | 1 |
| 7 | Construct CFG L= $\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} ; \mathrm{n} \geq 1\right\}$. | 3 | 2,3 | 6 |
| 8 | Find a LM derivation for aaabbabbba with the productions. | 3 | 2 | 5 |
| 9 | Find L(G), S $\rightarrow$ aSb, $\mathrm{S} \rightarrow \mathrm{ab}$. | 3 | 2 | 5 |
| 10 | Show that id* id can be generated by two distinct leftmost derivation in the grammar | 3 | 2 | 2 |


| 11 | Write a CFG for the set of strings which does not produce any palindromes. | 3 | 2 | 6 |
| :---: | :---: | :---: | :---: | :---: |
| 12 | Find the derivation tree for the grammar $\mathrm{G}=(\{\mathrm{S}, \mathrm{A}, \mathrm{B}\},\{\mathrm{a}, \mathrm{b}\}, \mathrm{P}, \mathrm{S}\}$ Where P is given by $\begin{aligned} & \mathrm{S} \rightarrow \mathrm{Aa} / \mathrm{bB} \\ & \mathrm{~A} \rightarrow \mathrm{ab} \\ & \mathrm{~B} \rightarrow \mathrm{aBb} / \mathrm{a} \\ & \hline \end{aligned}$ | 3 | 2 | 5 |
| 13 | Define parse tree. | 3 | 1 | 1 |
| 14 | What are the two major normal forms for context-free grammar? | 3 | 2 | 4 |
| 15 | What is a useless symbol? | 3 | 2 | 4 |
| 16 | Define Nullable Variable? | 3 | 1 | 1 |
| 17 | $\begin{aligned} & \text { Let } \mathrm{G}=(\mathrm{V}, \mathrm{~T}, \mathrm{P}, \mathrm{~S}) \text { with the productions given by } \\ & \mathrm{S} \rightarrow \mathrm{aSbS} / \mathrm{B} / \varepsilon \\ & \mathrm{B} \rightarrow \mathrm{abB} \\ & \text { Eliminate the useless production. } \\ & \hline \end{aligned}$ | 3 | 2 | 5 |
| 18 | What is a useful production? | 3 | 2 | 4 |
| 19 | Determine whether the grammar $G$ has a useless production? $\begin{aligned} & \mathrm{S} \rightarrow \mathrm{~A} \\ & \mathrm{~A} \rightarrow \mathrm{aA} / \varepsilon \\ & \mathrm{B} \rightarrow \mathrm{bA} \end{aligned}$ | 3 | 2 | 4 |
| 20 | Write a procedure to eliminate $\varepsilon$ production. | 3 | 2 | 6 |
| 21 | Write the procedure to eliminate the unit productions. | 3 | 2 | 6 |
| 22 | Define CNF. | 3 | 1 | 1 |
| 23 | Define GNF. | 3 | 1 | 1 |
| Part - B |  |  |  |  |
| 24 | Consider the Grammar G whose productions are $\begin{aligned} & \mathrm{S} \rightarrow 0 \mathrm{~B} / 1 \mathrm{~A} \\ & \mathrm{~A} \rightarrow 0 / 0 \mathrm{~S} / 1 \mathrm{AA} \end{aligned}$ <br> $\mathrm{B} \rightarrow 1 / 1 \mathrm{~S} / 0 \mathrm{BB}$ and the string 0110 <br> a. Find the left most derivation and associated derivation tree. <br> b. Find the right most derivation and associated derivation tree. <br> c. Show that the G is ambiguous. <br> d. Find L(G) | 3 | 2 | 5 |
| 25 | Consider the Grammar whose productions are $\begin{aligned} & \mathrm{S} \rightarrow \mathrm{aAS} / \mathrm{a} \\ & \mathrm{~A} \rightarrow \mathrm{SbA} / \mathrm{SS} / \mathrm{ba} \end{aligned}$ <br> a. Construct a LMD and RMD Tree for $\mathrm{S}=$ >* $^{*}$ aabbaa <br> b. Find the above grammar is ambiguous or unambiguous. | 3 | 2 | 5 |
| 26 | Construct Right Linear Grammar from the given Finite Automata | 3 | 2,3 | 6 |


|  | Construct Left Linear Grammar from the given Finite Automata <br> 27 |  |  |  |
| :--- | :--- | :--- | :--- | :--- |

## Unit - IV

| $\begin{aligned} & \hline \text { Sl. } \\ & \text { No } \end{aligned}$ | Questions | CO | PO | BT |
| :---: | :---: | :---: | :---: | :---: |
| Part - A |  |  |  |  |
| 1 | Define pushdown automaton. | 4 | 1 | 1 |
| 2 | What are the different ways of language acceptances by a PDA and define them. | 4 | 2 | 4 |
| 3 | Construct a PDA that accepts the language generated by the grammar $\mathrm{S} \rightarrow \mathrm{aSbb} / \mathrm{aab}$ | 4 | 2,3 | 6 |
| 4 | Construct a PDA that accepts the language generated by the grammar $\mathrm{S} \rightarrow \mathrm{aABB}, \mathrm{~A} \rightarrow \mathrm{aB} / \mathrm{a}, \mathrm{~B} \rightarrow \mathrm{bA} / \mathrm{b}$ | 4 | 2,3 | 6 |
| 5 | How do you convert CFG to a PDA. | 4 | 2 | 6 |
| 6 | Define Deterministic PDA. | 4 | 1 | 1 |
| 7 | Is it true that NDPA is more powerful than that od DPDA? Justify your answer. | 4 | 2 | 5 |
| 8 | Is it true that the language accepted by a PDA by empty stack and final states are different languages. | 4 | 2 | 5 |
| 9 | What is the additional feature PDA has when compared with NFA? Is PDA superior over NFA in the sense L acceptance? Justify your answer. | 4 | 2 | 4 |
| Part - B |  |  |  |  |
| 10 | Prove that if $\mathrm{L}=\mathrm{N}(\mathrm{PN})$ for some $\mathrm{PDA} \mathrm{PN}=(\mathrm{Q}, \Sigma, \Gamma, \delta, \mathrm{q} 0, \mathrm{Z} 0, \mathrm{~F})$, then there is a PDA PF such that $\mathrm{L}=\mathrm{L}(\mathrm{PF})$. | 4 | 1,2 | 2 |
| 11 | Prove that if $\mathrm{M} 1=(\mathrm{Q}, \Sigma, \Gamma, \delta, \mathrm{q} 0, \mathrm{Z} 0, \mathrm{~F})$ accept by final state, we can find a PDA M2, accepting $L$ by empty store i.e., $L=L(M 1)=N(M 2)$. | 4 | 1,2 | 2 |
| 12 | Construct a PDA that accepts the following languages <br> a. $\mathrm{L}=\left\{\mathrm{wcwr} \mid \mathrm{W}\right.$ in $\left.(0+1)^{*}\right\}$ by empty stack or final state <br> b. $L=\left\{w w R ; w \in(0+1)^{*}\right\}$ by empty stack or final state <br> c. $\mathrm{L}=\{0 \mathrm{n} 1 \mathrm{n} ; \mathrm{n} \geq 0\}$ accepted by empty stack or final state <br> d. $\mathrm{L}=\{$ anbmemdn; $\mathrm{n}, \mathrm{m} \geq 1\}$ accepted by empty store and check whether the string $\mathrm{w}=$ aaabcddd is accept or not. | 4 | 2,3 | 6 |
| 13 | Construct a PDA that will accept the language generated by the grammar $\mathrm{G}=(\{\mathrm{S}, \mathrm{A}\},\{\mathrm{a}, \mathrm{b}\}, \mathrm{P}, \mathrm{S})$ with the productions $\mathrm{S} \rightarrow \mathrm{AA} / \mathrm{a}, \mathrm{A} \rightarrow \mathrm{SA} / \mathrm{b}$ and test whether "abbabb" is in N(M). | 4 | 2,3 | 6 |
| 14 | Consider the grammar $\mathrm{G}=(\mathrm{V}, \mathrm{T}, \mathrm{P}, \mathrm{S})$ and test whether " 0101001 " is in $\mathrm{N}(\mathrm{M})$. Where Productions are $\begin{aligned} & \mathrm{S} \rightarrow 0 \mathrm{~S} / 1 \mathrm{~A} / 1 / 0 \mathrm{~B} / 0 \\ & \mathrm{~A} \rightarrow 0 \mathrm{~A} / 1 \mathrm{~B} / 0 / 1 \\ & \mathrm{~B} \rightarrow 0 \mathrm{~B} / 1 \mathrm{~A} / 0 / 1 \end{aligned}$ | 4 | 2,3 | 4 |
| 15 | Construct a PDA from the given $\mathrm{CFG} \mathrm{G}=(\{\mathrm{S}, \mathrm{A}\},\{\mathrm{a}, \mathrm{b}\}, \mathrm{P}, \mathrm{S})$ where the productions are $\mathrm{S} \rightarrow \mathrm{AS} / \varepsilon$ and $\mathrm{A} \rightarrow \mathrm{aAb} / \mathrm{Sb} / \mathrm{a}$ | 4 | 2,3 | 6 |
| 16 | Construct a PDA from the following CFG. <br> $\mathrm{G}=(\mathrm{V}, \mathrm{T}, \mathrm{P}, \mathrm{S})$ with $\mathrm{V}=\{\mathrm{S}\}, \mathrm{T}=\{\mathrm{a}, \mathrm{b}, \mathrm{c}\}$, and $\mathrm{P}=\{\mathrm{S} \rightarrow \mathrm{aSa}$, <br> $\mathrm{S} \rightarrow \mathrm{bSb}, \mathrm{S} \rightarrow \mathrm{c}\}$ | 4 | 2,3 | 6 |
| 17 | Convert the PDA $\mathrm{P}=(\{\mathrm{p}, \mathrm{q}\},\{0,1\},\{\mathrm{X}, \mathrm{Z} 0\}, \delta, \mathrm{q}, \mathrm{Z} 0)$ to a CFG . Where $\delta$ is given below: $\begin{array}{ll} \delta(\mathrm{q} 0,0, \mathrm{~S})=\{(\mathrm{q} 0, \mathrm{AS})\} & \delta(\mathrm{q} 0,0, \mathrm{~A})=\{(\mathrm{q} 0, \mathrm{AA}),(\mathrm{q} 1, \mathrm{~S})\} \\ \delta(\mathrm{q} 0,1, \mathrm{~A})=\{(\mathrm{q} 1, \varepsilon)\} & \delta(\mathrm{q} 1,1, \mathrm{~A})=\{(\mathrm{q} 1, \varepsilon)\} \\ \delta(\mathrm{q} 1, \varepsilon, \mathrm{~A})=\{(\mathrm{q} 1, \varepsilon)\} & \delta(\mathrm{q} 1, \varepsilon, \mathrm{~S})=\{(\mathrm{q} 1, \varepsilon)\} \end{array}$ | 4 | 2,3 | 5 |


|  | Convert the PDA P $=(\{\mathrm{p}, \mathrm{q}\},\{0,1\},\{\mathrm{X}, \mathrm{Z} 0\}, \delta, \mathrm{q}, \mathrm{Z} 0)$ <br> given by <br>  <br>  <br> $\delta(\mathrm{q}, 1, \mathrm{Z} 0)=\{(\mathrm{q}, \mathrm{XZ} 0)\}$ <br>  <br> $\delta(\mathrm{q}, 1, \mathrm{X})=\{(\mathrm{q}, \mathrm{XX})\}$ <br>  <br> $\delta(\mathrm{q}, 0, \mathrm{X})=\{(\mathrm{p}, \mathrm{X})\}$ <br>  <br> $\delta(\mathrm{q}, \varepsilon, \mathrm{X})=\{(\mathrm{q}, \varepsilon)\}$ <br> $\delta(\mathrm{p}, 1, \mathrm{X})=\{\{\mathrm{p}, \varepsilon)\}$ <br>  <br> $\delta(\mathrm{p}, 0, \mathrm{Z} 0)=\{(\mathrm{q}, \mathrm{Z} 0)\}$ |  |  |  |
| :--- | :--- | :--- | :--- | :--- |

## Unit - V

| $\begin{aligned} & \hline \text { Sl. } \\ & \text { No } \end{aligned}$ | Questions | CO | PO | BT |
| :---: | :---: | :---: | :---: | :---: |
| Part - A |  |  |  |  |
| 1 | What is a Turning Machine? | 5 | 1 | 4 |
| 2 | Define a Turing Machine. | 5 | 1 | 1 |
| 3 | Define Instantaneous description of TM. | 5 | 1 | 1 |
| 4 | What are the applications of TM? | 5 | 1 | 4 |
| 5 | What are the required fields of an instantaneous description or configuration of a TM. | 5 | 1 | 4 |
| 6 | Differentiate PDA and TM. | 5 | 2 | 3 |
| 7 | Define Universal TM | 5 | 1 | 1 |
| 8 | When is a function f said to be Turing computable? | 5 | 2 | 4 |
| 9 | Explain the Class of Grammars. | 5 | 1 | 2 |
| 10 | Discuss about PCP. | 5 | 2 | 2 |
| 11 | Differentiate PCP and MPCP. | 5 | 2 | 4 |
| Part - B |  |  |  |  |
| 12 | Design a TM to recognize the language $\mathrm{L}=\left\{\mathrm{a}^{\mathrm{n}} \mathrm{b}^{\mathrm{n}} ; \mathrm{n}>0\right\}$ and test whether the strings "aabb" is accepts or not. | 5 | 2,3 | 6 |
| 13 | Design a TM to recognize the language $\mathrm{L}=\left\{\mathrm{ww}^{\mathrm{r}} ; \mathrm{w} \in(\mathrm{a}+\mathrm{b})^{*}\right\}$ and test whether the strings "abba" is accepts or not. | 5 | 2,3 | 6 |
| 14 | Design a TM to recognize the language $\mathrm{L}=\left\{\mathrm{wcw}^{\mathrm{r}} ; \mathrm{w} \in(0+1)^{*}\right\}$. | 5 | 2,3 | 6 |
| 15 | Design a Turing machine to compute proper subtraction m-n. | 5 | 2,3 | 6 |
| 16 | Explain the class of Grammars with example. | 5 | 1,2 | 2 |
| 17 | Explain the PCP and MPCP with example | 5 | 1,2 | 2 |

